## Abstract simplicial complex

## Simplicial complex

A simplicial complex $\mathcal{C}$ is a set of finite sets such that if $\sigma \in \mathcal{C}$ and $\tau \subseteq \sigma$ then $\tau \in \mathcal{C}$. For every $\tau \subseteq \sigma \in \mathcal{C}$, the set $\tau$ is a face of $\sigma$, whereas $\sigma$ is a coface of $\tau$.

## k-simplex

$\sigma \in \mathcal{C}$ with $|\sigma|=k+1$ is called a $k$-simplex.

## Orientation

An orientation of $k$-simplices is an equivalence class of orderings where two simplices are considered equal if the permutation has a sign of 1 .

## Geometric realization

■ Realize a $k$-simplex as the convex hull of $k+1$ affinely independent points in some $\mathbb{R}^{d}$, with $d \geq k$.
■ Need to ensure that the simplicess only intersect along shared faces.

- Geometric intuition:

0 -simplices vertices
1-simplices edges
2-simplices triangles
3-simplices tetrahedra

Not interested in that.

## Chain group

## $k$ th chain group

The $k$ th chain group $C_{k}$ of $\mathcal{C}$ is the free abelian group on the set of oriented $k$-simplices. The group contains all abstract combinations of oriented $k$-simplices with coefficients from either a field or a principal ideal domain.
$c \in C_{k}$ is a $k$-chain, i.e.

$$
c=\sum_{i} \lambda_{i}\left[\sigma_{i}\right]
$$

with $\lambda_{i} \in \mathbb{Z}$, for example, and $\sigma_{i} \in \mathcal{C}$.

## Boundary operator

The $k$ th boundary operator $\partial_{k}: C_{k} \rightarrow C_{k-1}$ is a homomorphism whose action on a chain $c$ is defined on a simplex $\sigma=\left[v_{0}, v_{1}, \ldots, v_{k}\right]$ by

$$
\partial_{k} \sigma=\sum_{i}(-1)^{i}\left[v_{0}, v_{1}, \ldots, \hat{v}_{i}, \ldots, v_{k}\right]
$$

where $\hat{v}_{i}$ signifies that the $i$ th vertex is removed from the chain.

## Chain complex and subgroups

The boundary operators connect the chain groups of different dimensions. This forms a chain complex, i.e.

$$
\cdots \rightarrow C_{k+1} \rightarrow C_{k} \rightarrow C_{k-1} \rightarrow \ldots
$$

## Subgroups of $C_{k}$

We have the cycle group $Z_{k}=\operatorname{ker} \partial_{k}$ (mnemonic: "Zykel") and the boundary group $B_{k}=\operatorname{im} \partial_{k+1}$. Since $\partial_{k} \partial_{k+1}=0$, the subgroups are nested:

$$
B_{k} \subseteq Z_{k} \subseteq C_{k}
$$

## Homology

## $k$ th homology group

$$
H_{k}=Z_{k} / B_{k}
$$

This is well-defined because the subgroups are nested. The elements of the $k$ th homology group are classes of homologous cycles. If the coefficients are taken from a field $\mathbb{F}$ then $H_{k}$ becomes a vector space.

## Betti numbers

$$
\beta_{k}=\operatorname{rank} H_{k}
$$

- $\beta_{0}$ is the number of connected components
- $\beta_{1}$ is the number of 2-dimensional holes (circles)
- $\beta_{2}$ is the number of 3-dimensional holes (voids)


## Homology is useful

- Invariants of topological spaces

■ "Homology googles" to distinguish different spaces from one another

## Classical example

|  | Möbius strip | Torus |
| :--- | :--- | :--- |
| $H_{0}$ | $\mathbb{Z}$ | $\mathbb{Z}$ |
| $H_{1}$ | $\mathbb{Z}$ | $\mathbb{Z} \times \mathbb{Z}$ |
| $H_{2}$ | 0 | $\mathbb{Z}$ |



## Input data

## Assumption

The input data is given as a high-dimensional point cloud. There is some kind of metric, i.e. Euclidean distance.

## Goal

Identify "interesting" topological structures in the data-especially relevant for time series data.

## How to obtain a simplicial complex?

- Use points in point cloud as vertics of a graph

■ Determine edges by proximity, i.e. take all vertices situated within a distance of $\epsilon$

- This yields the neighbourhood graph $N_{\epsilon}$

■ Expand the graph afterwards

## Chěch complex

Topologically faithful but very hard to compute. Relies on precise distances.

## Vietoris-Rips complex

Less expensive calculation but possibly different homotopy type, i.e. we may not "see" what we want to see.

## How to choose $\epsilon$ ?



Figure: $\epsilon=0.013$


Figure: $\epsilon=0.019$

## Persistent homology

■ Need to distinguish between "essential" and "non-essential" holes

- Question of "optimal" values for $\epsilon$ is a mistake


## Idea

Do all computations for a large range of parameter values for $\epsilon$. Features that persist over the course of varying the parameter are likely to be "real" topological features.

## Visualization



Figure: Default "barcode" visualization taken from [1].

## Workflow (so far)



## Current status of my work

- Literature survey; we need to know the state of the art

■ Implemented algorithms for constructing the Vietoris-Rips complex [2]

- Started working on implementation of persistent homology calculation [3]


## Problems

- Complexes are very large
- Calculations are slow
- Not many applications out there (this may be a good thing)


## Roadmap

■ Even more literature survey
■ Examination of some data sets-how can we profit from these methods?

■ Try approximations to topology (sometimes we know the topology of the underlying space)
■ Rather vague: Use domain knowledge

## Possible applications

- Time-series data
- Clustered data

■?

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